

# Coloring Curves That Cross a Fixed Curve\*

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## Abstract

We prove that for every integer  $t \geq 1$ , the class of intersection graphs of curves in the plane each of which crosses a fixed curve in at least one and at most  $t$  points is  $\chi$ -bounded. This is essentially the strongest  $\chi$ -boundedness result one can get for this kind of graph classes. As a corollary, we prove that for any fixed integers  $k \geq 2$  and  $t \geq 1$ , every  $k$ -quasi-planar topological graph on  $n$  vertices with any two edges crossing at most  $t$  times has  $O(n \log n)$  edges.

**1998 ACM Subject Classification** I.3.5 Computational Geometry and Object Modeling

**Keywords and phrases** String graphs,  $\chi$ -boundedness,  $k$ -quasi-planar graphs

**Digital Object Identifier** 10.4230/LIPIcs.SoCG.2017.56

## 1 Introduction

### Overview

A *curve* is a homeomorphic image of the real interval  $[0, 1]$  in the plane. The *intersection graph* of a family of curves has these curves as vertices and the intersecting pairs of curves as edges. Combinatorial and algorithmic aspects of intersection graphs of curves, known as *string graphs*, have been attracting researchers for decades. A significant part of this research has been devoted to understanding classes of string graphs that are  $\chi$ -bounded, which means that every graph  $G$  in the class satisfies  $\chi(G) \leq f(\omega(G))$  for some function  $f: \mathbb{N} \rightarrow \mathbb{N}$ , where  $\chi(G)$  and  $\omega(G)$  denote the chromatic number and the clique number (the maximum size of a clique) of  $G$ , respectively. Recently, Pawlik et al. [24, 25] proved that the class of all string graphs is not  $\chi$ -bounded. However, all known constructions of string graphs with small clique number and large chromatic number require a lot of freedom in placing curves around in the plane.

What restrictions on placement of curves lead to  $\chi$ -bounded classes of intersection graphs? McGuinness [19, 20] proposed studying families of curves that cross a fixed curve *exactly once*. This initiated a series of results culminating in the proof that the class of intersection graphs of such families is indeed  $\chi$ -bounded [26]. By contrast, the class of intersection graphs of curves each crossing a fixed curve *at least once* is equal to the class of all string graphs and therefore is not  $\chi$ -bounded. We prove an essentially farthest possible generalization of the former result, allowing curves to cross the fixed curve *at least once and at most  $t$  times*, for any bound  $t$ .

► **Theorem 1.** *For every integer  $t \geq 1$ , the class of intersection graphs of curves each crossing a fixed curve in at least one and at most  $t$  points is  $\chi$ -bounded.*

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\* Alexandre Rok was partially supported by Israel Science Foundation grant 1136/12. Bartosz Walczak was partially supported by National Science Center of Poland grant 2015/17/D/ST1/00585.



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33rd International Symposium on Computational Geometry (SoCG 2017).

Editors: Boris Aronov and Matthew J. Katz; Article No. 56; pp. 56:1–56:15

Leibniz International Proceedings in Informatics



LIPICs Schloss Dagstuhl – Leibniz-Zentrum für Informatik, Dagstuhl Publishing, Germany

Additional motivation for Theorem 1 comes from its application to bounding the number of edges in so-called  $k$ -quasi-planar graphs, which we discuss at the end of this introduction.

### Context

Chromatic number of intersection graphs of geometric objects has been investigated since the 1960s. In a seminal paper, Asplund and Grünbaum [3] proved that intersection graphs of axis-parallel rectangles in the plane satisfy  $\chi = O(\omega^2)$  and conjectured that for every integer  $d \geq 1$ , there is a function  $f_d: \mathbb{N} \rightarrow \mathbb{N}$  such that intersection graphs of axis-parallel boxes in  $\mathbb{R}^d$  satisfy  $\chi \leq f_d(\omega)$ . However, a few years later, a surprising construction due to Burling [5] showed that there are triangle-free intersection graphs of axis-parallel boxes in  $\mathbb{R}^3$  with arbitrarily large chromatic number. Since then, the upper bound of  $O(\omega^2)$  and the trivial lower bound of  $\Omega(\omega)$  on the maximum possible chromatic number of a rectangle intersection graph have been improved only in terms of multiplicative constants [11, 13].

Another classical example of a  $\chi$ -bounded class of geometric intersection graphs is provided by circle graphs—intersection graphs of chords of a fixed circle. Gyárfás [10] proved that circle graphs satisfy  $\chi = O(\omega^2 4^\omega)$ . The best known upper and lower bounds on the maximum possible chromatic number of a circle graph are  $O(2^\omega)$  [14] and  $\Omega(\omega \log \omega)$  [12].

McGuinness [19, 20] proposed investigating the problem when much more general geometric shapes are allowed but the way how they are arranged in the plane is restricted. In [19], he proved that the class of intersection graphs of L-shapes crossing a fixed horizontal line is  $\chi$ -bounded. Families of L-shapes in the plane are *simple*, which means that any two members of the family intersect in at most one point. McGuinness [20] also showed that triangle-free intersection graphs of simple families of curves each crossing a fixed line in exactly one point have bounded chromatic number. Further progress in this direction was made by Suk [27], who proved that simple families of  $x$ -monotone curves crossing a fixed vertical line give rise to a  $\chi$ -bounded class of intersection graphs, and by Lasoń et al. [17], who reached the same conclusion without assuming that the curves are  $x$ -monotone. Finally, in [26], we proved that the class of intersection graphs of curves each crossing a fixed line in exactly one point is  $\chi$ -bounded. These results remain valid if the fixed straight line is replaced by a fixed curve [28].

The class of string graphs is not  $\chi$ -bounded. Pawlik et al. [24, 25] presented a construction of triangle-free intersection graphs of segments (or geometric shapes of various other kinds) with chromatic number growing as fast as  $\Theta(\log \log n)$  with the number of vertices  $n$ . It was further generalized to a construction of string graphs with clique number  $\omega$  and chromatic number  $\Theta_\omega((\log \log n)^{\omega-1})$  [16]. The best known upper bound on the chromatic number of string graphs in terms of the number of vertices is  $(\log n)^{O(\log \omega)}$ , proved by Fox and Pach [8] using a separator theorem for string graphs due to Matoušek [18]. For intersection graphs of segments or, more generally,  $x$ -monotone curves, an upper bound of the form  $\chi = O_\omega(\log n)$  follows from the above-mentioned result in [27] or [26] via recursive halving. Upper bounds of the form  $\chi = O_\omega((\log \log n)^{f(\omega)})$  (for some function  $f: \mathbb{N} \rightarrow \mathbb{N}$ ) are known for very special classes of string graphs: rectangle overlap graphs [15, 16] and subtree overlap graphs [16]. The former still allow the triangle-free construction with  $\chi = \Theta(\log \log n)$  and the latter the construction with  $\chi = \Theta_\omega((\log \log n)^{\omega-1})$ .

### Quasi-planarity

A *topological graph* is a graph with a fixed curvilinear drawing in the plane. For  $k \geq 2$ , a  *$k$ -quasi-planar graph* is a topological graph with no  $k$  pairwise crossing edges. In particular, a 2-quasi-planar graph is just a planar graph. It is conjectured that  $k$ -quasi-planar graphs with

$n$  vertices have  $O_k(n)$  edges [4, 23]. For  $k = 2$ , this asserts a well-known property of planar graphs. The conjecture is also verified for  $k = 3$  [2, 22] and  $k = 4$  [1], but it remains open for  $k \geq 5$ . Best known upper bounds on the number of edges in a  $k$ -quasi-planar graph are  $n(\log n)^{O(\log k)}$  in general [7, 8],  $O_k(n \log n)$  for the case of  $x$ -monotone edges [29],  $O_k(n \log n)$  for the case that any two edges intersect at most once [28], and  $2^{\alpha(n)^\nu} n \log n$  for the case that any two edges intersect in at most  $t$  points, where  $\alpha$  is the inverse Ackermann function and  $\nu$  depends on  $k$  and  $t$  [28]. We apply Theorem 1 to improve the last bound to  $O_{k,t}(n \log n)$ .

► **Theorem 2.** *Every  $k$ -quasi-planar topological graph  $G$  on  $n$  vertices such that any two edges of  $G$  intersect in at most  $t$  points has at most  $\mu_{k,t} n \log n$  edges, where  $\mu_{k,t}$  depends only on  $k$  and  $t$ .*

The proof follows the same line as the proof in [28] for the case  $t = 1$  (see Section 3).

## 2 Proof of Theorem 1

### Setup

Let  $\mathbb{N}$  denote the set of positive integers. Graph-theoretic terms applied to a family of curves  $\mathcal{F}$  have the same meaning as applied to the intersection graph of  $\mathcal{F}$ . In particular, the *chromatic number* of  $\mathcal{F}$ , denoted by  $\chi(\mathcal{F})$ , is the minimum number of colors in a *proper coloring* of  $\mathcal{F}$  (a coloring that distinguishes pairs of intersecting curves), and the *clique number* of  $\mathcal{F}$ , denoted by  $\omega(\mathcal{F})$ , is the maximum size of a *clique* in  $\mathcal{F}$  (a set of pairwise intersecting curves in  $\mathcal{F}$ ).

► **Theorem 1 (rephrased).** *For every  $t \in \mathbb{N}$ , there is a non-decreasing function  $f_t: \mathbb{N} \rightarrow \mathbb{N}$  with the following property: for any fixed curve  $c_0$ , every family  $\mathcal{F}$  of curves each intersecting  $c_0$  in at least one and at most  $t$  points satisfies  $\chi(\mathcal{F}) \leq f_t(\omega(\mathcal{F}))$ .*

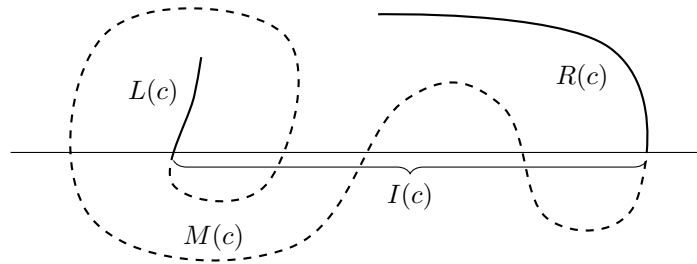
A point  $p$  is a *proper crossing* of curves  $c_1$  and  $c_2$  if  $c_1$  passes from one side to the other side of  $c_2$  in a sufficiently small neighborhood of  $p$ . From now on, without significant loss of generality, we make the following implicit assumption: any two distinct curves that we consider intersect in finitely many points, and each of their intersection points is a proper crossing. There is one exception to the latter condition: a curve  $c$  may have an endpoint on another curve if this is required by the definition of  $c$  (like for 1-curves defined below).

### Initial reduction

We start by reducing Theorem 1 to a somewhat simpler and more convenient setting. We fix a horizontal line in the plane and call it the *baseline*. The upper half-plane bounded by the baseline is denoted by  $H^+$ . A *1-curve* is a curve in  $H^+$  that has one endpoint on the baseline and does not intersect the baseline in any other point. Intersection graphs of 1-curves are known as *outerstring graphs* and form a  $\chi$ -bounded class of graphs—this result, due to the authors, is the starting point of the proof of Theorem 1.

► **Theorem 3 ([26]).** *There is a non-decreasing function  $f_0: \mathbb{N} \rightarrow \mathbb{N}$  such that every family  $\mathcal{F}$  of 1-curves satisfies  $\chi(\mathcal{F}) \leq f_0(\omega(\mathcal{F}))$ .*

An *even-curve* is a curve that has both endpoints above the baseline and intersects the baseline in at least two points (this is an even number, by the proper crossing assumption). For  $t \in \mathbb{N}$ , a *2t-curve* is an even-curve that intersects the baseline in exactly  $2t$  points. The *basepoint* of a 1-curve  $s$  is the endpoint of  $s$  on the baseline. A *basepoint* of an even-curve  $c$



■ **Figure 1**  $L(c)$ ,  $R(c)$ ,  $M(c)$  (all the dashed part), and  $I(c)$  for a 6-curve  $c$ .

is an intersection point of  $c$  with the baseline. Every even-curve  $c$  determines two 1-curves—the two parts of  $c$  from an endpoint to the closest basepoint. They are called the 1-curves of  $c$  and denoted by  $L(c)$  and  $R(c)$  so that the basepoint of  $L(c)$  lies to the left of the basepoint of  $R(c)$  on the baseline (see Figure 1). A family  $\mathcal{F}$  of even-curves is an *LR-family* if every intersection between two curves  $c_1, c_2 \in \mathcal{F}$  is an intersection between  $L(c_1)$  and  $R(c_2)$  or between  $L(c_2)$  and  $R(c_1)$ . The main effort in this paper goes to proving the following statement on *LR-families* of even-curves.

► **Theorem 4.** *There is a non-decreasing function  $f: \mathbb{N} \rightarrow \mathbb{N}$  such that every LR-family  $\mathcal{F}$  of even-curves satisfies  $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F}))$ .*

Theorem 4 makes no assumption on the maximum number of intersection points of an even-curve with the baseline. We derive Theorem 1 from Theorem 4 in two steps, first proving the following lemma, and then showing that Theorem 1 is essentially a special case of it.

► **Lemma 5.** *For every  $t \in \mathbb{N}$ , there is a non-decreasing function  $f_t: \mathbb{N} \rightarrow \mathbb{N}$  such that every family  $\mathcal{F}$  of  $2t$ -curves no two of which intersect on or below the baseline satisfies  $\chi(\mathcal{F}) \leq f_t(\omega(\mathcal{F}))$ .*

**Proof of Lemma 5 from Theorem 4.** The proof goes by induction on  $t$ . Let  $f_0$  and  $f$  be the functions claimed by Theorem 3 and Theorem 4, respectively, and let  $f_t(k) = f_{t-1}^2(k)f(k)$  for  $t \geq 1$  and  $k \in \mathbb{N}$ . We establish the base case for  $t = 1$  and the induction step for  $t \geq 2$  simultaneously. Namely, fix an integer  $t \geq 1$ , and let  $\mathcal{F}$  be as in the statement of the lemma. For every  $2t$ -curve  $c \in \mathcal{F}$ , enumerate the endpoints and basepoints of  $c$  as  $p_0(c), \dots, p_{2t+1}(c)$  in their order along  $c$  so that  $p_0(c)$  and  $p_1(c)$  are the endpoints of  $L(c)$  while  $p_{2t}(c)$  and  $p_{2t+1}(c)$  are the endpoints of  $R(c)$ . Build two families of curves  $\mathcal{F}_1$  and  $\mathcal{F}_2$  putting the part of  $c$  from  $p_0(c)$  to  $p_{2t-1}(c)$  to  $\mathcal{F}_1$  and the part of  $c$  from  $p_2(c)$  to  $p_{2t+1}(c)$  to  $\mathcal{F}_2$  for every  $c \in \mathcal{F}$ . If  $t = 1$ , then  $\mathcal{F}_1$  and  $\mathcal{F}_2$  are families of 1-curves. If  $t \geq 2$ , then  $\mathcal{F}_1$  and  $\mathcal{F}_2$  are equivalent to families of  $2(t-1)$ -curves, because the curve in  $\mathcal{F}_1$  or  $\mathcal{F}_2$  obtained from a  $2t$ -curve  $c \in \mathcal{F}$  can be shortened a little at  $p_{2t-1}(c)$  or  $p_2(c)$ , respectively, losing that basepoint but no intersection points with other curves. Therefore, by Theorem 3 or the induction hypothesis, we have  $\chi(\mathcal{F}_k) \leq f_{t-1}(\omega(\mathcal{F}_k)) \leq f_{t-1}(\omega(\mathcal{F}))$  for  $k \in \{1, 2\}$ . For  $c \in \mathcal{F}$  and  $k \in \{1, 2\}$ , let  $\phi_k(c)$  be the color of the curve obtained from  $c$  in an optimal proper coloring of  $\mathcal{F}_k$ . Every subfamily of  $\mathcal{F}$  on which  $\phi_1$  and  $\phi_2$  are constant is an *LR-family* and therefore, by Theorem 4 and monotonicity of  $f$ , has chromatic number at most  $f(\omega(\mathcal{F}))$ . We conclude that  $\chi(\mathcal{F}) \leq \chi(\mathcal{F}_1)\chi(\mathcal{F}_2)f(\omega(\mathcal{F})) \leq f_{t-1}^2(\omega(\mathcal{F}))f(\omega(\mathcal{F})) = f_t(\omega(\mathcal{F}))$ . ◀

A *closed curve* is a homeomorphic image of a unit circle in the plane. For a closed curve  $\gamma$ , the Jordan curve theorem asserts that the set  $\mathbb{R}^2 \setminus \gamma$  consists of two connected components: one bounded, denoted by  $\text{int } \gamma$ , and one unbounded, denoted by  $\text{ext } \gamma$ .

**Proof of Theorem 1 from Theorem 4.** We elect to present this proof in an intuitive rather than rigorous way. Let  $\mathcal{F}$  be a family of curves each intersecting  $c_0$  in at least one and at most  $t$  points. Let  $\gamma_0$  be a closed curve surrounding  $c_0$  very closely so that  $\gamma_0$  intersects every curve in  $\mathcal{F}$  in exactly  $2t$  points (winding if necessary to increase the number of intersections) and all endpoints of curves in  $\mathcal{F}$  and intersection points of pairs of curves in  $\mathcal{F}$  lie in  $\text{ext } \gamma_0$ . We “invert”  $\text{int } \gamma_0$  with  $\text{ext } \gamma_0$  to obtain an equivalent family of curves  $\mathcal{F}'$  and a closed curve  $\gamma'_0$  with the same properties except that all endpoints of curves in  $\mathcal{F}'$  and intersection points of pairs of curves in  $\mathcal{F}'$  lie in  $\text{int } \gamma'_0$ . It follows that some part of  $\gamma'_0$  lies in the unbounded component of  $\mathbb{R}^2 \setminus \bigcup \mathcal{F}'$ . We “cut”  $\gamma'_0$  there and “unfold” it into the baseline, transforming  $\mathcal{F}'$  into an equivalent family  $\mathcal{F}''$  of  $2t$ -curves all endpoints of which and intersection points of pairs of which lie above the baseline. The “equivalence” of  $\mathcal{F}$ ,  $\mathcal{F}'$ , and  $\mathcal{F}''$  means in particular that the intersection graphs of  $\mathcal{F}$ ,  $\mathcal{F}'$ , and  $\mathcal{F}''$  are isomorphic, so the theorem follows from Lemma 5 (and thus Theorem 4). ◀

A statement analogous to Theorem 4 fails for families of objects each consisting of two 1-curves only, without the “middle part” connecting them. Specifically, we define a *double-curve* as a set  $X \subset H^+$  that is a union of two disjoint 1-curves, denoted by  $L(X)$  and  $R(X)$  so that the basepoint of  $L(X)$  lies to the left of the basepoint of  $R(X)$ , and we call a family  $\mathcal{X}$  of double-curves an *LR-family* if every intersection between two double-curves  $X_1, X_2 \in \mathcal{X}$  is an intersection between  $L(X_1)$  and  $R(X_2)$  or between  $L(X_2)$  and  $R(X_1)$ .

► **Theorem 6.** *For every  $\zeta \in \mathbb{N}$ , there is a triangle-free LR-family of double-curves  $\mathcal{X}$  such that  $\chi(\mathcal{X}) \geq \zeta$ .*

The proof of Theorem 6 is an easy adaptation of the construction from [24, 25]. We omit the details. The rest of this section is devoted to the proof of Theorem 4.

### Overview of the proof of Theorem 4

Recall the assertion of Theorem 4: the *LR-families* of even-curves are  $\chi$ -bounded. The proof is quite long and technical, so we find it useful to provide a high-level overview of its structure. The proof will be presented via a series of reductions. First, we will reduce Theorem 4 to the following statement (Lemma 7): the *LR-families* of 2-curves are  $\chi$ -bounded. This statement will be proved by induction on the clique number. Specifically, we will prove the following as the induction step: if every *LR-family* of 2-curves  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k - 1$  satisfies  $\chi(\mathcal{F}) \leq \xi$ , then every *LR-family* of 2-curves  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$  satisfies  $\chi(\mathcal{F}) \leq \zeta$ , where  $\zeta$  is a constant depending only on  $k$  and  $\xi$ . The only purpose of the induction hypothesis is to infer that if  $\omega(\mathcal{F}) \leq k$  and  $c \in \mathcal{F}$ , then the family of 2-curves in  $\mathcal{F} \setminus \{c\}$  that intersect  $c$  has chromatic number at most  $\xi$ . For notational convenience, *LR-families* of 2-curves with the latter property will be called  $\xi$ -families. We will thus reduce the problem to the following statement (Lemma 9): the  $\xi$ -families are  $\chi$ -bounded, where the  $\chi$ -bounding function depends on  $\xi$ .

We will deal with  $\xi$ -families via a series of technical lemmas of the following general form: every  $\xi$ -family with chromatic number large enough contains a specific configuration of curves. Two kinds of such configurations are particularly important: (a) a large clique, and (b) a 2-curve  $c$  and a subfamily  $\mathcal{F}'$  with large chromatic number such that the basepoints of the 2-curves in  $\mathcal{F}'$  lie between the basepoints of  $c$ . In the core of the argument are the proofs that

- every  $\xi$ -family with chromatic number large enough contains (a) or (b) (Lemma 16),
- assuming the above, every  $\xi$ -family with chromatic number large enough contains (a).

Combined, they complete the argument. Since the two proofs are almost identical, we introduce one more reduction—to  $(\xi, h)$ -families (Lemma 15). A  $(\xi, h)$ -family is just a  $\xi$ -family that satisfies an additional technical condition sufficient to carry both proofs at once.

**More notation and terminology**

Let  $\prec$  denote the left-to-right order of points on the baseline ( $p_1 \prec p_2$  means that  $p_1$  is to the left of  $p_2$ ). For convenience, we also use the notation  $\prec$  for curves intersecting the baseline ( $c_1 \prec c_2$  means that every basepoint of  $c_1$  is to the left of every basepoint of  $c_2$ ) and for families of such curves ( $\mathcal{C}_1 \prec \mathcal{C}_2$  means that  $c_1 \prec c_2$  for any  $c_1 \in \mathcal{C}_1$  and  $c_2 \in \mathcal{C}_2$ ). For a family  $\mathcal{C}$  of curves intersecting the baseline (even-curves or 1-curves) and two 1-curves  $x$  and  $y$ , let  $\mathcal{C}(x, y) = \{c \in \mathcal{C} : x \prec c \prec y\}$  or  $\mathcal{C}(x, y) = \{c \in \mathcal{C} : y \prec c \prec x\}$  depending on whether  $x \prec y$  or  $y \prec x$ . For a family  $\mathcal{C}$  of curves intersecting the baseline and a segment  $I$  on the baseline, let  $\mathcal{C}(I)$  denote the family of curves in  $\mathcal{C}$  with all basepoints on  $I$ .

For an even-curve  $c$ , let  $M(c)$  denote the subcurve of  $c$  connecting the basepoints of  $L(c)$  and  $R(c)$ , and let  $I(c)$  denote the segment on the baseline connecting the basepoints of  $L(c)$  and  $R(c)$  (see Figure 1). For a family  $\mathcal{F}$  of even-curves, let  $L(\mathcal{F}) = \{L(c) : c \in \mathcal{F}\}$ ,  $R(\mathcal{F}) = \{R(c) : c \in \mathcal{F}\}$ , and  $I(\mathcal{F})$  denote the minimal segment on the baseline that contains  $I(c)$  for every  $c \in \mathcal{F}$ .

A *cap-curve* is a curve in  $H^+$  that has both endpoints on the baseline and does not intersect the baseline in any other point. For a cap-curve  $\gamma$ , it follows from the Jordan curve theorem that the set  $H^+ \setminus \gamma$  consists of two connected components: one bounded, denoted by  $\text{int } \gamma$ , and one unbounded, denoted by  $\text{ext } \gamma$ . Any two cap-curves one with endpoints  $p_1, q_1$  and the other with endpoints  $p_2, q_2$  such that  $p_1 \prec p_2 \prec q_1 \prec q_2$  intersect in an odd number of points.

**Reduction to LR-families of 2-curves**

We will reduce Theorem 4 to the following statement on LR-families of 2-curves, which is essentially a special case of Theorem 4.

► **Lemma 7.** *There is a non-decreasing function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that every LR-family  $\mathcal{F}$  of 2-curves satisfies  $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F}))$ .*

A *component* of a family of 1-curves  $\mathcal{S}$  is a connected component of  $\bigcup \mathcal{S}$  (the union of all curves in  $\mathcal{S}$ ). The following easy but powerful observation reuses an idea from [17, 20, 27].

► **Lemma 8.** *For every LR-family of even-curves  $\mathcal{F}$ , if  $\mathcal{F}^*$  is the family of curves  $c \in \mathcal{F}$  such that  $L(c)$  and  $R(c)$  lie in distinct components of  $L(\mathcal{F}) \cup R(\mathcal{F})$ , then  $\chi(\mathcal{F}^*) \leq 4$ .*

**Proof.** Let  $G$  be an auxiliary graph where the vertices are the components of  $L(\mathcal{F}) \cup R(\mathcal{F})$  and the edges are the pairs  $V_1V_2$  of components such that there is a curve  $c \in \mathcal{F}^*$  with  $L(c) \subseteq V_1$  and  $R(c) \subseteq V_2$  or  $L(c) \subseteq V_2$  and  $R(c) \subseteq V_1$ . Since  $\mathcal{F}$  is an LR-family, the curves in  $\mathcal{F}^*$  cannot intersect “outside” the components of  $L(\mathcal{F}) \cup R(\mathcal{F})$ . It follows that  $G$  is planar and thus 4-colorable. Fix a proper 4-coloring of  $G$ , and assign the color of a component  $V$  to every curve  $c \in \mathcal{F}^*$  with  $L(c) \subseteq V$ . For any  $c_1, c_2 \in \mathcal{F}^*$ , if  $L(c_1)$  and  $R(c_2)$  intersect, then  $L(c_1)$  and  $R(c_2)$  lie in the same component  $V_1$  while  $L(c_2)$  lies in a component  $V_2$  such that  $V_1V_2$  is an edge of  $G$ , so  $c_1$  and  $c_2$  are assigned distinct colors. The coloring of  $\mathcal{F}^*$  is therefore proper. ◀

**Proof of Theorem 4 from Lemma 7.** We show that  $\chi(\mathcal{F}) \leq f(\omega(\mathcal{F})) + 4$ , where  $f$  is the function claimed by Lemma 7. We have  $\mathcal{F} = \mathcal{F}_1 \cup \mathcal{F}_2$ , where  $\mathcal{F}_1 = \{c \in \mathcal{F} : L(c) \text{ and } R(c) \text{ lie in the same component of } L(\mathcal{F}) \cup R(\mathcal{F})\}$  and  $\mathcal{F}_2 = \{c \in \mathcal{F} : L(c) \text{ and } R(c) \text{ lie in distinct components of } L(\mathcal{F}) \cup R(\mathcal{F})\}$ . Lemma 8 yields  $\chi(\mathcal{F}_2) \leq 4$ . It remains to show that  $\chi(\mathcal{F}_1) \leq f(\omega(\mathcal{F}))$ .

Let  $c_1, c_2 \in \mathcal{F}_1$ . We claim that the intervals  $I(c_1)$  and  $I(c_2)$  are nested or disjoint. Suppose they are not. For  $\varepsilon > 0$  and a component  $V$  of  $L(\mathcal{F}) \cup R(\mathcal{F})$ , let  $V^\varepsilon$  denote the  $\varepsilon$ -neighborhood of  $V$  in  $H^+$ . We assume that  $\varepsilon$  is small enough so that the sets  $V^\varepsilon$  for all

components  $V$  of  $L(\mathcal{F}) \cup R(\mathcal{F})$  and the curves  $M(c)$  for all  $c \in \mathcal{F}_1$  are pairwise disjoint (except at common basepoints). For  $k \in \{1, 2\}$ , since  $L(c_k)$  and  $R(c_k)$  belong to the same component  $V_k$  of  $L(\mathcal{F}) \cup R(\mathcal{F})$ , there is a cap-curve  $\gamma_k \subset V_k^\varepsilon$  that connects the basepoints of  $L(c_k)$  and  $R(c_k)$ . We can assume without loss of generality that  $\gamma_1$  and  $\gamma_2$  intersect in a finite number of points and each of their intersection points is a proper crossing (this is why we take  $\gamma_k \subset V_k^\varepsilon$  instead of  $\gamma_k \subseteq V_k$ ). Since  $I(c_1)$  and  $I(c_2)$  are neither nested nor disjoint, the basepoints of  $L(c_2)$  and  $R(c_2)$  lie one in  $\text{int } \gamma_1$  and the other in  $\text{ext } \gamma_1$ , so  $\gamma_1$  and  $\gamma_2$  intersect in an odd number of points. For  $k \in \{1, 2\}$ , let  $\tilde{\gamma}_k$  be the closed curve obtained as the union of  $\gamma_k$  and  $M(c_k)$ . It follows that  $\tilde{\gamma}_1$  and  $\tilde{\gamma}_2$  intersect in an odd number of points and each of their intersection points is a proper crossing, which is a contradiction.

Transform  $\mathcal{F}_1$  into a family of 2-curves  $\mathcal{F}'_1$  replacing the part  $M(c)$  of every curve  $c \in \mathcal{F}_1$  by the lower semicircle connecting the endpoints of  $M(c)$ . These semicircles are pairwise disjoint (because  $I(c_1)$  and  $I(c_2)$  are nested or disjoint for any  $c_1, c_2 \in \mathcal{F}_1$ ), so  $\mathcal{F}'_1$  is an  $LR$ -family with intersection graph isomorphic to that of  $\mathcal{F}_1$ . Lemma 7 yields  $\chi(\mathcal{F}_1) = \chi(\mathcal{F}'_1) \leq f(\omega(\mathcal{F}'_1)) \leq f(\omega(\mathcal{F}))$ . ◀

### Reduction to $\xi$ -families

For  $\xi \in \mathbb{N}$ , a  $\xi$ -family is an  $LR$ -family of 2-curves  $\mathcal{F}$  with the following property: for every 2-curve  $c \in \mathcal{F}$ , the family of 2-curves in  $\mathcal{F} \setminus \{c\}$  that intersect  $c$  has chromatic number at most  $\xi$ . We reduce Lemma 7 to the following statement on  $\xi$ -families.

► **Lemma 9.** *For any  $\xi, k \in \mathbb{N}$ , there is a constant  $\zeta \in \mathbb{N}$  such that every  $\xi$ -family  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$  satisfies  $\chi(\mathcal{F}) \leq \zeta$ .*

**Proof of Lemma 7 from Lemma 9.** Let  $f(1) = 1$ . For  $k \geq 2$ , let  $f(k)$  be the constant claimed by Lemma 9 such that every  $f(k-1)$ -family  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$  satisfies  $\chi(\mathcal{F}) \leq f(k)$ . Let  $k = \omega(\mathcal{F})$ , and proceed by induction on  $k$  to prove  $\chi(\mathcal{F}) \leq f(k)$ . Clearly, if  $k = 1$ , then  $\chi(\mathcal{F}) = 1$ . For the induction step, assume  $k \geq 2$ . For every  $c \in \mathcal{F}$ , the family of 2-curves in  $\mathcal{F} \setminus \{c\}$  that intersect  $c$  has clique number at most  $k-1$  and therefore, by the induction hypothesis, has chromatic number at most  $f(k-1)$ . That is,  $\mathcal{F}$  is an  $f(k-1)$ -family, and the definition of  $f$  yields  $\chi(\mathcal{F}) \leq f(k)$ . ◀

### Dealing with $\xi$ -families

First, we establish the following special case of Lemma 9.

► **Lemma 10.** *For every  $\xi \in \mathbb{N}$ , every  $\xi$ -family  $\mathcal{F}$  with  $\bigcap_{c \in \mathcal{F}} I(c) \neq \emptyset$  satisfies  $\chi(\mathcal{F}) \leq 4\xi + 4$ .*

The proof of Lemma 10 is essentially the same as the proof of Lemma 19 in [28]. We need the following elementary lemma, which was also used in various forms in [17, 19, 20, 26, 27].

► **Lemma 11** (McGuinness [19, Lemma 2.1]). *Let  $G$  be a graph,  $\prec$  be a total order on the vertices of  $G$ , and  $\alpha, \beta \in \mathbb{N}$ . If  $\chi(G) > (2\beta + 2)\alpha$ , then  $G$  has an induced subgraph  $H$  such that  $\chi(H) > \alpha$  and  $\chi(G(u, v)) > \beta$  for every edge  $uv$  of  $H$ . In particular, if  $\chi(G) > 2\beta + 2$ , then  $G$  has an edge  $uv$  with  $\chi(G(u, v)) > \beta$ . Here,  $G(u, v)$  denotes the subgraph of  $G$  induced on the vertices strictly between  $u$  and  $v$  in the order  $\prec$ .*

**Proof of Lemma 10.** Suppose  $\chi(\mathcal{F}) > 4\xi + 4$ . Since  $\bigcap_{c \in \mathcal{F}} I(c) \neq \emptyset$ , the 2-curves in  $\mathcal{F}$  can be enumerated as  $c_1, \dots, c_{|\mathcal{F}|}$  so that  $L(c_1) \prec \dots \prec L(c_{|\mathcal{F}|}) \prec R(c_{|\mathcal{F}|}) \prec \dots \prec R(c_1)$ . Apply Lemma 11 to the intersection graph of  $\mathcal{F}$  and the order  $c_1, \dots, c_{|\mathcal{F}|}$  to obtain two indices  $i, j \in \{1, \dots, |\mathcal{F}|\}$  such that the 2-curves  $c_i$  and  $c_j$  intersect and  $\chi(\{c_{i+1}, \dots, c_{j-1}\}) > 2\xi + 1$ .



Assume  $L(c_i)$  and  $R(c_j)$  intersect; the argument for the other case is analogous. There is a cap-curve  $\gamma \subseteq L(c_i) \cup R(c_j)$  connecting the basepoints of  $L(c_i)$  and  $R(c_j)$ . Every curve intersecting  $\gamma$  intersects  $c_i$  or  $c_j$ . Since  $\mathcal{F}$  is a  $\xi$ -family, the 2-curves in  $\{c_{i+1}, \dots, c_{j-1}\}$  that intersect  $c_i$  have chromatic number at most  $\xi$ , and so do those that intersect  $c_j$ . Every 2-curve  $c_k \in \{c_{i+1}, \dots, c_{j-1}\}$  not intersecting  $\gamma$  satisfies  $L(c_k) \subset \text{int } \gamma$  and  $R(c_k) \subset \text{ext } \gamma$ , so these 2-curves are pairwise disjoint. We conclude that  $\chi(\{c_{i+1}, \dots, c_{j-1}\}) \leq 2\xi + 1$ , which is a contradiction.  $\blacktriangleleft$

Lemma 11 easily implies that every family of 2-curves  $\mathcal{F}$  with  $\chi(\mathcal{F}) > (2\beta + 2)^2\alpha$  contains a subfamily  $\mathcal{H}$  with  $\chi(\mathcal{H}) > \alpha$  such that  $\chi(\mathcal{F}(L(c_1), L(c_2))) > \beta$  and  $\chi(\mathcal{F}(R(c_1), R(c_2))) > \beta$  for any two intersecting 2-curves  $c_1, c_2 \in \mathcal{H}$ . This is considerably strengthened by the following lemma. Its proof extends the idea used in [19] for the proof of Lemma 11.

**► Lemma 12.** *For every  $\xi \in \mathbb{N}$ , there is a function  $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  with the following property: for any  $\alpha, \beta \in \mathbb{N}$  and every  $\xi$ -family  $\mathcal{F}$  with  $\chi(\mathcal{F}) > f(\alpha, \beta)$ , there is a subfamily  $\mathcal{H} \subseteq \mathcal{F}$  such that  $\chi(\mathcal{H}) > \alpha$  and  $\chi(\mathcal{F}(x, y)) > \beta$  for any two intersecting 1-curves  $x \in R(\mathcal{H})$  and  $y \in L(\mathcal{H})$ .*

**Proof.** Let  $f(\alpha, \beta) = (2\beta + 12\xi + 20)\alpha$ . Let  $\mathcal{F}$  be a  $\xi$ -family with  $\chi(\mathcal{F}) > f(\alpha, \beta)$ . Construct a sequence of points  $p_0 \prec \dots \prec p_{m+1}$  on the baseline with the following properties:

- the points  $p_0, \dots, p_{m+1}$  are distinct from all basepoints of 2-curves in  $\mathcal{F}$ ,
- $p_0$  lies to the left of and  $p_{m+1}$  lies to the right of all basepoints of 2-curves in  $\mathcal{F}$ ,
- $\chi(\mathcal{F}(p_i p_{i+1})) = \beta + 1$  for  $0 \leq i \leq m - 1$ , and  $\chi(\mathcal{F}(p_m p_{m+1})) \leq \beta + 1$ .

This is done greedily by first choosing  $p_1$  so that  $\chi(\mathcal{F}(p_0 p_1)) = \beta + 1$ , then choosing  $p_2$  so that  $\chi(\mathcal{F}(p_1 p_2)) = \beta + 1$ , and so on. For  $0 \leq i \leq j \leq m$ , let  $\mathcal{F}_{i,j} = \{c \in \mathcal{F} : p_i \prec L(c) \prec p_{i+1} \text{ and } p_j \prec R(c) \prec p_{j+1}\}$ . In particular,  $\mathcal{F}_{i,i} = \mathcal{F}(p_i p_{i+1})$  for  $0 \leq i \leq m$ . Since  $\mathcal{F} = \bigcup_{0 \leq i \leq j \leq m} \mathcal{F}_{i,j}$ , at least one of the following holds:

$$\chi(\bigcup_{i=0}^m \mathcal{F}_{i,i}) > (2\beta + 2)\alpha, \quad \chi(\bigcup_{i=0}^{m-1} \mathcal{F}_{i,i+1}) > (12\xi + 12)\alpha, \quad \chi(\bigcup_{i=0}^{m-2} \bigcup_{j=i+2}^m \mathcal{F}_{i,j}) > 6\alpha.$$

In each case, we will find a subfamily  $\mathcal{H} \subseteq \mathcal{F}$  such that any two intersecting 1-curves  $x \in R(\mathcal{H})$  and  $y \in L(\mathcal{H})$  satisfy  $x \in R(\mathcal{F}_{i,j})$  and  $y \in L(\mathcal{F}_{r,s})$ , where  $0 \leq i \leq j \leq m$ ,  $0 \leq r \leq s \leq m$ , and  $|j - r| \geq 2$ . Then,  $\chi(\mathcal{F}(x, y)) \geq \chi(\mathcal{F}(p_{\max(j,r)-1} p_{\max(j,r)})) = \beta + 1$ , as required.

Suppose  $\chi(\bigcup_{i=0}^m \mathcal{F}_{i,i}) > (2\beta + 2)\alpha$ . We have  $\chi(\mathcal{F}_{i,i}) \leq \beta + 1$  for  $0 \leq i \leq m$ . Color the 2-curves in every  $\mathcal{F}_{i,i}$  properly using the same set of  $\beta + 1$  colors on  $\mathcal{F}_{i,i}$  and  $\mathcal{F}_{r,r}$  whenever  $i \equiv r \pmod{2}$ , thus using  $2\beta + 2$  colors in total. It follows that  $\chi(\mathcal{H}) > \alpha$  for some family  $\mathcal{H} \subseteq \bigcup_{i=0}^m \mathcal{F}_{i,i}$  of 2-curves of the same color. To conclude, for any two intersecting 1-curves  $x \in R(\mathcal{H})$  and  $y \in L(\mathcal{H})$ , we have  $x \in R(\mathcal{F}_{i,i})$  and  $y \in L(\mathcal{F}_{r,r})$  for some distinct indices  $i, r \in \{0, \dots, m\}$  with  $i \equiv r \pmod{2}$  and thus  $|i - r| \geq 2$ .

Now, suppose  $\chi(\bigcup_{i=0}^{m-1} \mathcal{F}_{i,i+1}) > (12\xi + 12)\alpha$ . By Lemma 10, we have  $\chi(\mathcal{F}_{i,i+1}) \leq 4\xi + 4$  for  $0 \leq i \leq m - 1$ . Color the 2-curves in every  $\mathcal{F}_{i,i+1}$  properly using the same set of  $4\xi + 4$  colors on  $\mathcal{F}_{i,i+1}$  and  $\mathcal{F}_{r,r+1}$  whenever  $i \equiv r \pmod{3}$ , thus using  $12\xi + 12$  colors in total. It follows that  $\chi(\mathcal{H}) > \alpha$  for some family  $\mathcal{H} \subseteq \bigcup_{i=0}^{m-1} \mathcal{F}_{i,i+1}$  of 2-curves of the same color. To conclude, for any two intersecting 1-curves  $x \in R(\mathcal{H})$  and  $y \in L(\mathcal{H})$ , we have  $x \in R(\mathcal{F}_{i,i+1})$  and  $y \in L(\mathcal{F}_{r,r+1})$  for some distinct indices  $i, r \in \{0, \dots, m - 1\}$  with  $i \equiv r \pmod{3}$  and thus  $|i + 1 - r| \geq 2$ .

Finally, suppose  $\chi(\bigcup_{i=0}^{m-2} \bigcup_{j=i+2}^m \mathcal{F}_{i,j}) > 6\alpha$ . It follows that  $\chi(\bigcup_{i \in I} \bigcup_{j=i+2}^m \mathcal{F}_{i,j}) > 3\alpha$ , where  $I = \{i \in \{0, \dots, m - 2\} : i \equiv 0 \pmod{2}\}$  or  $I = \{i \in \{0, \dots, m - 2\} : i \equiv 1 \pmod{2}\}$ . Consider an auxiliary graph  $G$  with vertex set  $I$  and edge set  $\{ij : i, j \in I, i < j, \text{ and } \mathcal{F}_{i,j-1} \cup \mathcal{F}_{i,j} \neq \emptyset\}$ . Since no two 2-curves in  $\mathcal{F}$  cross below the baseline,  $G$  has no two edges  $i_1 j_1$  and  $i_2 j_2$  such that  $i_1 < i_2 < j_1 < j_2$ . In particular,  $G$  is an outerplanar graph, and



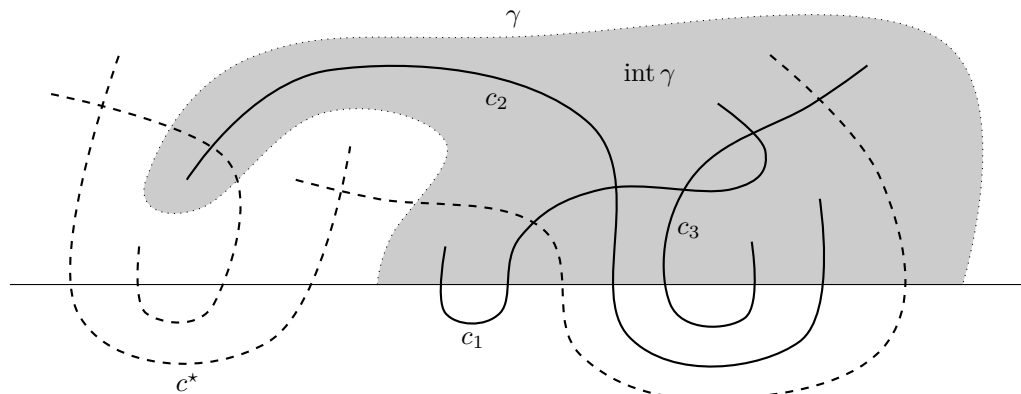


Figure 2 Illustration for Lemma 14:  $\mathcal{G} = \{c_1, c_2, c_3\}$ .

thus  $\chi(G) \leq 3$ . Fix a proper 3-coloring of  $G$ , and use the color of  $i$  on every 2-curve in  $\bigcup_{j=i+2}^m \mathcal{F}_{i,j}$  for every  $i \in I$ . It follows that  $\chi(\mathcal{H}) > \alpha$  for some family  $\mathcal{H} \subseteq \bigcup_{i \in I} \bigcup_{j=i+2}^m \mathcal{F}_{i,j}$  of 2-curves of the same color. To conclude, for any two intersecting 1-curves  $x \in R(\mathcal{H})$  and  $y \in L(\mathcal{H})$ , we have  $x \in R(\mathcal{F}_{i,j})$  and  $y \in L(\mathcal{F}_{r,s})$  for some indices  $i, r \in I, j \in \{i+2, \dots, m\}$ , and  $s \in \{r+2, \dots, m\}$  such that  $j \notin \{r-1, r\}$  (otherwise  $ir$  would be an edge of  $G$ ),  $j \neq r+1$  (otherwise two 2-curves, one from  $\mathcal{F}_{i,r+1}$  and one from  $\mathcal{F}_{r,s}$ , would cross below the baseline), and thus  $|j-r| \geq 2$ . ◀

It is proved in [26] that for every family of 1-curves  $\mathcal{S}$ , there are a cap-curve  $\gamma$  and a subfamily  $\mathcal{U} \subseteq \mathcal{S}$  with  $\chi(\mathcal{U}) \geq \frac{1}{2}\chi(\mathcal{S})$  such that every 1-curve in  $\mathcal{U}$  is contained in  $\text{int } \gamma$  and intersects some 1-curve in  $\mathcal{S}$  that intersects  $\text{ext } \gamma$ . The proof follows an idea from [10], used subsequently also in [17, 19, 20, 21, 27], where  $\mathcal{U}$  is chosen as one of the sets of 1-curves at a fixed distance from an appropriately chosen 1-curve in the intersection graph of  $\mathcal{S}$ , and  $\gamma$  is a cap-curve surrounding  $\mathcal{U}$  very closely. However, this method fails to imply an analogous statement for 2-curves. We will need a more powerful tool—part of the recent series of works on induced subgraphs that must be present in graphs with sufficiently large chromatic number.

► **Theorem 13** (Chudnovsky, Scott, Seymour [6, Theorem 1.8]). *There is a function  $f: \mathbb{N} \rightarrow \mathbb{N}$  with the following property: for every  $\alpha \in \mathbb{N}$ , every string graph  $G$  with  $\chi(G) > f(\alpha)$  contains a vertex  $v$  such that  $\chi(G_v^2) > \alpha$ , where  $G_v^2$  denotes the subgraph of  $G$  induced on the vertices within distance at most 2 from  $v$ .*

The special case of Theorem 13 for triangle-free intersection graphs of curves any two of which intersect in at most one point was proved earlier by McGuinness [21, Theorem 5.3].

► **Lemma 14** (see Figure 2). *For every  $\xi \in \mathbb{N}$ , there is a function  $f: \mathbb{N} \rightarrow \mathbb{N}$  with the following property: for every  $\alpha \in \mathbb{N}$  and every  $\xi$ -family  $\mathcal{F}$  with  $\chi(\mathcal{F}) > f(\alpha)$ , there are a cap-curve  $\gamma$  and a subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > \alpha$  such that every 2-curve  $c \in \mathcal{G}$  satisfies  $L(c), R(c) \subset \text{int } \gamma$  and intersects some 2-curve in  $\mathcal{F}$  that intersects  $\text{ext } \gamma$ .*

**Proof.** Let  $f(\alpha) = f_1(3\alpha + 5\xi + 5)$ , where  $f_1$  is the function claimed by Theorem 13. Let  $\mathcal{F}$  be a  $\xi$ -family with  $\chi(\mathcal{F}) > f(\alpha)$ . It follows that there is a 2-curve  $c^* \in \mathcal{F}$  such that the family of curves within distance at most 2 from  $c^*$  in the intersection graph of  $\mathcal{F}$  has chromatic number greater than  $3\alpha + 5\xi + 5$ . For  $k \in \{1, 2\}$ , let  $\mathcal{F}_k$  be the 2-curves in  $\mathcal{F}$  at distance exactly  $k$  from  $c^*$  in the intersection graph of  $\mathcal{F}$ . Since  $\chi(\{c^*\} \cup \mathcal{F}_1 \cup \mathcal{F}_2) > 3\alpha + 5\xi + 5$  and  $\chi(\mathcal{F}_1) \leq \xi$  (because  $\mathcal{F}$  is a  $\xi$ -family), we have  $\chi(\mathcal{F}_2) > 3\alpha + 4\xi + 4$ . We have  $\mathcal{F}_2 = \mathcal{G}_1 \cup \mathcal{G}_2 \cup \mathcal{G}_3 \cup \mathcal{G}_4$ , where

$\mathcal{G}_1 = \{c \in \mathcal{F}_2 : L(c) \prec R(c) \prec L(c^*) \prec R(c^*)\}$ ,  $\mathcal{G}_2 = \{c \in \mathcal{F}_2 : L(c^*) \prec L(c) \prec R(c) \prec R(c^*)\}$ ,  
 $\mathcal{G}_3 = \{c \in \mathcal{F}_2 : L(c^*) \prec R(c^*) \prec L(c) \prec R(c)\}$ ,  $\mathcal{G}_4 = \{c \in \mathcal{F}_2 : L(c) \prec L(c^*) \prec R(c^*) \prec R(c)\}$ .  
 Since  $\chi(\mathcal{F}_2) > 3\alpha + 4\xi + 4$  and  $\chi(\mathcal{G}_4) \leq 4\xi + 4$  (by Lemma 10), we have  $\chi(\mathcal{G}_k) > \alpha$  for  
 some  $k \in \{1, 2, 3\}$ . Since neither basepoint of  $c^*$  lies on  $I(\mathcal{G}_k)$ , there is a cap-curve  $\gamma$  with  
 $L(c^*), R(c^*) \subset \text{ext } \gamma$  and  $L(c), R(c) \subset \text{int } \gamma$  for all  $c \in \mathcal{G}_k$ . The lemma follows with  $\mathcal{G} = \mathcal{G}_k$ . ◀

### Reduction to $(\xi, h)$ -families

For  $\xi \in \mathbb{N}$  and a function  $h: \mathbb{N} \rightarrow \mathbb{N}$ , a  $(\xi, h)$ -family is a  $\xi$ -family  $\mathcal{F}$  with the following  
 additional property: for every  $\alpha \in \mathbb{N}$  and every subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > h(\alpha)$ , there is  
 a subfamily  $\mathcal{H} \subseteq \mathcal{G}$  with  $\chi(\mathcal{H}) > \alpha$  such that every 2-curve in  $\mathcal{F}$  with a basepoint on  $I(\mathcal{H})$   
 has both basepoints on  $I(\mathcal{G})$ . We will prove the following lemma.

► **Lemma 15.** *For any  $\xi, k \in \mathbb{N}$  and any function  $h: \mathbb{N} \rightarrow \mathbb{N}$ , there is a constant  $\zeta \in \mathbb{N}$  such  
 that every  $(\xi, h)$ -family  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$  satisfies  $\chi(\mathcal{F}) \leq \zeta$ .*

The notion of a  $(\xi, h)$ -family and Lemma 15 provide a convenient abstraction of what is  
 needed to prove the next lemma and then to prove Lemma 9 with the use of the next lemma.

► **Lemma 16.** *For any  $\xi, k \in \mathbb{N}$ , there is a function  $f: \mathbb{N} \rightarrow \mathbb{N}$  such that for every  $\alpha \in \mathbb{N}$ ,  
 every  $\xi$ -family  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$  and  $\chi(\mathcal{F}) > f(\alpha)$  contains a 2-curve  $c$  with  $\chi(\mathcal{F}(I(c))) > \alpha$ .*

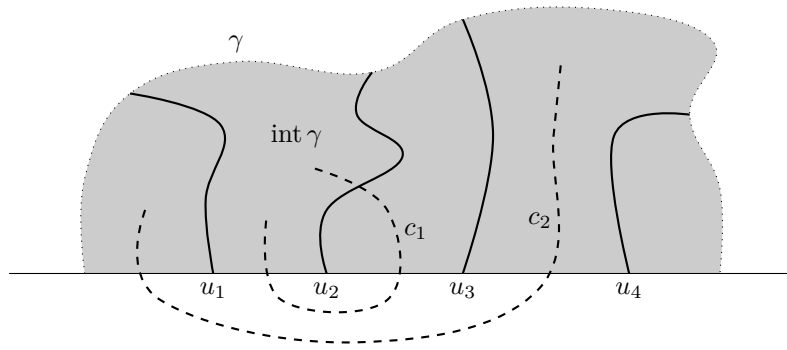
**Proof of Lemma 16 from Lemma 15.** Let  $h_\alpha: \mathbb{N} \ni \beta \mapsto \beta + 2\alpha + 2 \in \mathbb{N}$ , and let  $f(\alpha)$   
 be the constant claimed by Lemma 15 such that every  $(\xi, h_\alpha)$ -family  $\mathcal{F}$  with  $\omega(\mathcal{F}) \leq k$   
 satisfies  $\chi(\mathcal{F}) \leq f(\alpha)$ . Let  $\mathcal{F}$  be a  $\xi$ -family with  $\omega(\mathcal{F}) \leq k$  and  $\chi(\mathcal{F}(I(c))) \leq \alpha$  for every  
 $c \in \mathcal{F}$ . It is enough to show that  $\mathcal{F}$  is a  $(\xi, h_\alpha)$ -family. To this end, consider a subfamily  
 $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > h_\alpha(\beta)$  for some  $\beta \in \mathbb{N}$ . Take  $\mathcal{G}_L, \mathcal{G}_R \subseteq \mathcal{G}$  so that  $L(\mathcal{G}_L) \prec L(\mathcal{G} \setminus \mathcal{G}_L)$ ,  
 $\chi(\mathcal{G}_L) = \alpha + 1$ ,  $R(\mathcal{G} \setminus \mathcal{G}_R) \prec R(\mathcal{G}_R)$ , and  $\chi(\mathcal{G}_R) = \alpha + 1$ . Let  $\mathcal{H} = \mathcal{G} \setminus (\mathcal{G}_L \cup \mathcal{G}_R)$ . It follows  
 that  $\chi(\mathcal{H}) \geq \chi(\mathcal{G}) - 2\alpha - 2 > \beta$ . If there is a 2-curve  $c \in \mathcal{F}$  with one basepoint on  $I(\mathcal{H})$  and  
 the other basepoint not on  $I(\mathcal{G})$ , then  $\mathcal{G}_L \subseteq \mathcal{F}(I(c))$  or  $\mathcal{G}_R \subseteq \mathcal{F}(I(c))$ , so  $\chi(\mathcal{F}(I(c))) \geq \alpha + 1$ ,  
 which is a contradiction. Therefore, every 2-curve in  $\mathcal{F}$  with a basepoint on  $I(\mathcal{H})$  has both  
 basepoints on  $I(\mathcal{G})$ . This shows that  $\mathcal{F}$  is a  $(\xi, h_\alpha)$ -family. ◀

**Proof of Lemma 9 from Lemma 15.** Let  $h$  be the function claimed by Lemma 16 for  $\xi$  and  
 $k$ . Let  $\mathcal{F}$  be a  $\xi$ -family with  $\omega(\mathcal{F}) \leq k$ . In view of Lemma 15, it is enough to show that  
 $\mathcal{F}$  is a  $(\xi, h)$ -family. To this end, consider a subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > h(\alpha)$  for some  
 $\alpha \in \mathbb{N}$ . Lemma 16 yields a 2-curve  $c \in \mathcal{G}$  such that  $\chi(\mathcal{G}(I(c))) > \alpha$ . Every 2-curve in  $\mathcal{F}$   
 with a basepoint on  $I(c)$  has both basepoints on  $I(c)$ , otherwise it would cross  $c$  below the  
 baseline. Therefore, the condition of a  $(\xi, h)$ -family is satisfied with  $\mathcal{H} = \mathcal{G}(I(c))$ . ◀

### Dealing with $(\xi, h)$ -families

The rest of the proof is inspired from the ideas in [26]. A family of 1-curves  $\mathcal{S}$  *supports* a  
 family of 2-curves  $\mathcal{F}$  if every 2-curve in  $\mathcal{F}$  intersects some 1-curve in  $\mathcal{S}$ . A *skeleton* is a pair  
 $(\gamma, \mathcal{U})$  such that  $\gamma$  is a cap-curve and  $\mathcal{U}$  is a family of pairwise disjoint 1-curves each of which  
 has one endpoint (other than the basepoint) on  $\gamma$  and all the remaining part in  $\text{int } \gamma$  (see  
 Figure 3). For a family of 1-curves  $\mathcal{S}$ , a skeleton  $(\gamma, \mathcal{U})$  is an  $\mathcal{S}$ -skeleton if every 1-curve in  $\mathcal{U}$   
 is a subcurve of some 1-curve in  $\mathcal{S}$ . A skeleton  $(\gamma, \mathcal{U})$  *supports* a family of 2-curves  $\mathcal{F}$  if every  
 2-curve  $c \in \mathcal{F}$  satisfies  $L(c), R(c) \subset \text{int } \gamma$  and intersects some 1-curve in  $\mathcal{U}$ .

► **Lemma 17.** *For every function  $h: \mathbb{N} \rightarrow \mathbb{N}$ , there is a function  $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  such that  
 for any  $\alpha, \beta \in \mathbb{N}$ , every  $(\xi, h)$ -family  $\mathcal{F}$  with  $\chi(\mathcal{F}) > f(\alpha, \beta)$  contains one of the following  
 configurations:*



■ **Figure 3** A skeleton  $(\gamma, \{u_1, u_2, u_3, u_4\})$ , which supports  $c_1$  but not  $c_2$ .

- a subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > \alpha$  supported by an  $L(\mathcal{F})$ -skeleton or an  $R(\mathcal{F})$ -skeleton,
- a subfamily  $\mathcal{H} \subseteq \mathcal{F}$  with  $\chi(\mathcal{H}) > \beta$  supported by a family of 1-curves  $\mathcal{S}$  with  $\mathcal{S} \subseteq L(\mathcal{F})$  or  $\mathcal{S} \subseteq R(\mathcal{F})$  such that  $s \prec \mathcal{H}$  or  $\mathcal{H} \prec s$  for every 1-curve  $s \in \mathcal{S}$ .

**Proof.** Let  $f(\alpha, \beta) = f_1(2\alpha + h(2\beta) + 4)$ , where  $f_1$  is the function claimed by Lemma 14. Apply Lemma 14 to obtain a cap-curve  $\gamma$  and a subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > 2\alpha + h(2\beta) + 4$  such that every 2-curve  $c \in \mathcal{G}$  satisfies  $L(c), R(c) \subset \text{int } \gamma$  and intersects some 2-curve in  $\mathcal{F}_{\text{ext}}$ . Here and further on,  $\mathcal{F}_{\text{ext}}$  denotes the family of 2-curves in  $\mathcal{F}$  that intersect  $\text{ext } \gamma$ . Let  $\mathcal{U}_L$  be the 1-curves that are subcurves of 1-curves in  $L(\mathcal{F})$ , have one endpoint (other than the basepoint) on  $\gamma$ , and have all the remaining part in  $\text{int } \gamma$ . Let  $\mathcal{U}_R$  be the 1-curves that are subcurves of 1-curves in  $R(\mathcal{F})$ , have one endpoint (other than the basepoint) on  $\gamma$ , and have all the remaining part in  $\text{int } \gamma$ . Thus  $(\gamma, \mathcal{U}_L)$  is an  $L(\mathcal{F})$ -skeleton, and  $(\gamma, \mathcal{U}_R)$  is an  $R(\mathcal{F})$ -skeleton. Let  $\mathcal{G}_L$  be the 2-curves in  $\mathcal{G}$  that intersect some 1-curve in  $\mathcal{U}_L$ , and let  $\mathcal{G}_R$  be those that intersect some 1-curve in  $\mathcal{U}_R$ . If  $\chi(\mathcal{G}_L) > \alpha$  or  $\chi(\mathcal{G}_R) > \alpha$ , then the first conclusion of the lemma holds. Thus assume  $\chi(\mathcal{G}_L) \leq \alpha$  and  $\chi(\mathcal{G}_R) \leq \alpha$ . Let  $\mathcal{G}' = \mathcal{G} \setminus (\mathcal{G}_L \cup \mathcal{G}_R)$ . It follows that  $\chi(\mathcal{G}') \geq \chi(\mathcal{G}) - 2\alpha > h(2\beta) + 4$ .

By Lemma 8, the 2-curves  $c \in \mathcal{G}'$  such that  $L(c)$  and  $R(c)$  lie in distinct components of  $L(\mathcal{G}') \cup R(\mathcal{G}')$  have chromatic number at most 4. Therefore, there is a component  $V$  of  $L(\mathcal{G}') \cup R(\mathcal{G}')$  such that  $\chi(\mathcal{G}'_V) \geq \chi(\mathcal{G}') - 4 > h(2\beta)$ , where  $\mathcal{G}'_V = \{c \in \mathcal{G}' : L(c), R(c) \subseteq V\}$ . There is a cap-curve  $\nu \subseteq V$  connecting the two endpoints of the segment  $I(\mathcal{G}'_V)$ . Suppose there is a 2-curve  $c \in \mathcal{F}_{\text{ext}}$  with both basepoints on  $I(\mathcal{G}'_V)$ . If  $L(c)$  intersects  $\text{ext } \gamma$ , then the part of  $L(c)$  from the basepoint to the first intersection point with  $\gamma$ , which is a 1-curve in  $\mathcal{U}_L$ , must intersect  $\nu$  (as  $\nu \subseteq V \subset \text{int } \gamma$ ) and thus a curve in  $\mathcal{G}'$  (as  $V$  is a component of  $\mathcal{G}'$ ). Thus  $\mathcal{G}' \cap \mathcal{G}_L \neq \emptyset$ , which is a contradiction. An analogous contradiction is reached if  $R(c)$  intersects  $\text{ext } \gamma$ . This shows that no curve in  $\mathcal{F}_{\text{ext}}$  has both basepoints on  $I(\mathcal{G}'_V)$ .

Since  $\mathcal{F}$  is a  $(\xi, h)$ -family and  $\chi(\mathcal{G}'_V) > h(2\beta)$ , there is a subfamily  $\mathcal{H}' \subseteq \mathcal{G}'_V$  such that  $\chi(\mathcal{H}') > 2\beta$  and every 2-curve in  $\mathcal{F}$  with a basepoint on  $I(\mathcal{H}')$  has the other basepoint on  $I(\mathcal{G}'_V)$ . This and the above imply that no curve in  $\mathcal{F}_{\text{ext}}$  has a basepoint on  $I(\mathcal{H}')$ . Since every curve in  $\mathcal{H}'$  intersects some curve in  $\mathcal{F}_{\text{ext}}$ , we have  $\mathcal{H}' = \mathcal{H}_L \cup \mathcal{H}_R$ , where  $\mathcal{H}_L$  are the 2-curves in  $\mathcal{H}'$  that intersect some 1-curve in  $L(\mathcal{F}_{\text{ext}})$  and  $\mathcal{H}_R$  are those that intersect some 1-curve in  $R(\mathcal{F}_{\text{ext}})$ . Since  $\chi(\mathcal{H}') > 2\beta$ , we conclude that  $\chi(\mathcal{H}_L) > \beta$  or  $\chi(\mathcal{H}_R) > \beta$  and thus the second conclusion of the lemma holds with  $(\mathcal{H}, \mathcal{S}) = (\mathcal{H}_L, L(\mathcal{F}_{\text{ext}}))$  or  $(\mathcal{H}, \mathcal{S}) = (\mathcal{H}_R, R(\mathcal{F}_{\text{ext}}))$ . ◀

► **Lemma 18.** For every function  $h: \mathbb{N} \rightarrow \mathbb{N}$ , there is a function  $f: \mathbb{N} \rightarrow \mathbb{N}$  such that for every  $\alpha \in \mathbb{N}$ , every  $(\xi, h)$ -family  $\mathcal{F}$  with  $\chi(\mathcal{F}) > f(\alpha)$  contains a subfamily  $\mathcal{G} \subseteq \mathcal{F}$  with  $\chi(\mathcal{G}) > \alpha$  supported by an  $L(\mathcal{F})$ -skeleton or an  $R(\mathcal{F})$ -skeleton.

**Proof.** Let  $f(\alpha) = f_1(\alpha, f_1(\alpha, f_1(\alpha, 4\xi)))$ , where  $f_1$  is the function claimed by Lemma 17. Suppose to the contrary that no such subfamily  $\mathcal{G}$  exists. Let  $\mathcal{F}_0 = \mathcal{F}$ . Apply Lemma 17 three times to obtain families  $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3, \mathcal{S}_1, \mathcal{S}_2$ , and  $\mathcal{S}_3$  with the following properties:

- $\mathcal{F} = \mathcal{F}_0 \supseteq \mathcal{F}_1 \supseteq \mathcal{F}_2 \supseteq \mathcal{F}_3$ ,
- for  $i \in \{1, 2, 3\}$ , we have  $\mathcal{S}_i \subseteq L(\mathcal{F}_{i-1})$  or  $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$ ,  $\mathcal{F}_i$  is supported by  $\mathcal{S}_i$ , and  $s \prec \mathcal{F}_i$  or  $\mathcal{F}_i \prec s$  for every 1-curve  $s \in \mathcal{S}_i$ .
- $\chi(\mathcal{F}_1) > f_1(\alpha, f_1(\alpha, 4\xi))$ ,  $\chi(\mathcal{F}_2) > f_1(\alpha, 4\xi)$  and  $\chi(\mathcal{F}_3) > 4\xi$ .

There are indices  $i, j \in \{1, 2, 3\}$  with  $i < j$  such that  $\mathcal{S}_i$  and  $\mathcal{S}_j$  are of the same ‘‘type’’: either  $\mathcal{S}_i \subseteq L(\mathcal{F}_{i-1})$  and  $\mathcal{S}_j \subseteq L(\mathcal{F}_{j-1})$  or  $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$  and  $\mathcal{S}_j \subseteq R(\mathcal{F}_{j-1})$ . Assume for the rest of the proof that  $\mathcal{S}_i \subseteq R(\mathcal{F}_{i-1})$  and  $\mathcal{S}_j \subseteq R(\mathcal{F}_{j-1})$ ; the argument for the other case is analogous.

Let  $\mathcal{S}_< = \{s \in \mathcal{S}_j : s \prec \mathcal{F}_j\}$ ,  $\mathcal{S}_> = \{s \in \mathcal{S}_j : \mathcal{F}_j \prec s\}$ ,  $\mathcal{F}_<$  be the 2-curves in  $\mathcal{F}_j$  that intersect some 1-curve in  $\mathcal{S}_<$ , and  $\mathcal{F}_>$  be those that intersect some 1-curve in  $\mathcal{S}_>$ . Thus  $\mathcal{F}_< \cup \mathcal{F}_> = \mathcal{F}_j$ . This and  $\chi(\mathcal{F}_j) \geq \chi(\mathcal{F}_3) > 4\xi$  yield  $\chi(\mathcal{F}_<) > 2\xi$  or  $\chi(\mathcal{F}_>) > 2\xi$ . Assume for the rest of the proof that  $\chi(\mathcal{F}_<) > 2\xi$ ; the argument for the other case is analogous.

Let  $\mathcal{S}_<^{\min}$  be an inclusion-minimal subfamily of  $\mathcal{S}_<$  with the property that  $\mathcal{S}_<^{\min}$  still supports  $\mathcal{F}_<$ . Let  $s^*$  be the 1-curve in  $\mathcal{S}_<^{\min}$  with rightmost basepoint, and let  $\mathcal{F}_<^* = \{c \in \mathcal{F}_< : L(c) \text{ intersects } s^*\}$ . Since  $\mathcal{F}$  is a  $\xi$ -family, we have  $\chi(\mathcal{F}_<^*) \leq \xi$ . By the choice of  $\mathcal{S}_<^{\min}$ , there exists a 2-curve  $c^* \in \mathcal{F}_<^*$  disjoint from every 1-curve in  $\mathcal{S}_<^{\min}$  other than  $s^*$ . Since  $\mathcal{F}_<$  is supported by  $\mathcal{S}_i$ , there is a 1-curve  $s_i \in \mathcal{S}_i$  that intersects  $L(c^*)$ . We show that every 2-curve in  $\mathcal{F}_< \setminus \mathcal{F}_<^*$  intersects  $s_i$ .

Let  $c \in \mathcal{F}_< \setminus \mathcal{F}_<^*$ , and let  $s$  be a 1-curve in  $\mathcal{S}_<^{\min}$  that intersects  $L(c)$ . Thus  $s \neq s^*$ , by the definition of  $\mathcal{F}_<^*$ . There is a cap-curve  $\gamma \subseteq L(c) \cup s$ . Since  $s \prec s^* \prec L(c)$  and  $s^*$  intersects neither  $s$  nor  $L(c)$ , we have  $s^* \subset \text{int } \gamma$ . Since  $L(c^*)$  intersects  $s^*$  but neither  $s$  nor  $L(c)$ , we also have  $L(c^*) \subset \text{int } \gamma$ . Since  $s_i \prec \mathcal{F}_i$  or  $\mathcal{F}_i \prec s_i$ , the basepoint of  $s_i$  lies in  $\text{ext } \gamma$ . Therefore, since  $s_i$  intersects  $L(c^*)$ , the 1-curve  $s_i$  must enter  $\text{int } \gamma$  through a point on  $L(c)$ . This shows that every 2-curve in  $\mathcal{F}_< \setminus \mathcal{F}_<^*$  intersects  $s_i$ . This and the assumption that  $\mathcal{F}$  is a  $\xi$ -family yield  $\chi(\mathcal{F}_< \setminus \mathcal{F}_<^*) \leq \xi$ . We conclude that  $\chi(\mathcal{F}_<) \leq \chi(\mathcal{F}_<^*) + \chi(\mathcal{F}_< \setminus \mathcal{F}_<^*) \leq 2\xi$ , which is a contradiction. ◀

A *chain* of length  $n$  is a sequence  $((a_1, b_1), \dots, (a_n, b_n))$  of pairs of 2-curves such that

- for  $1 \leq i \leq n$ , the 1-curves  $R(a_i)$  and  $L(b_i)$  intersect,
- for  $2 \leq i \leq n$ , the basepoints of  $R(a_i)$  and  $L(b_i)$  lie between the basepoints of  $R(a_{i-1})$  and  $L(b_{i-1})$ , and  $L(a_i)$  intersects  $R(a_1), \dots, R(a_{i-1})$  or  $R(b_i)$  intersects  $L(b_1), \dots, L(b_{i-1})$ .

► **Lemma 19.** *For every  $\xi \in \mathbb{N}$  and every function  $h: \mathbb{N} \rightarrow \mathbb{N}$ , there is a function  $f: \mathbb{N} \rightarrow \mathbb{N}$  such that for every  $n \in \mathbb{N}$ , every  $(\xi, h)$ -family  $\mathcal{F}$  with  $\chi(\mathcal{F}) > f(n)$  contains a chain of length  $n$ .*

**Proof (see Figure 4).** We define the function  $f$  by induction. Let  $f(1) = 1$ ; if  $\chi(\mathcal{F}) > 1$ , then  $\mathcal{F}$  contains two intersecting 2-curves, which form a chain of length 1. For the induction step, fix  $n \geq 1$ , and assume that every  $(\xi, h)$ -family  $\mathcal{H}$  with  $\chi(\mathcal{H}) > f(n)$  contains a chain of length  $n$ . Let  $\beta = f_1(f(n), h(2\xi) + 4\xi + 2)$  and  $f(n+1) = f_2(f_2(f_2(\beta)))$ , where  $f_1$  is the function claimed by Lemma 12 and  $f_2$  is the function claimed by Lemma 18. Let  $\mathcal{F}$  be a  $(\xi, h)$ -family with  $\chi(\mathcal{F}) > f(n+1)$ . We claim that  $\mathcal{F}$  contains a chain of length  $n+1$ .

Let  $\mathcal{F}_0 = \mathcal{F}$ . Apply Lemma 18 three times to find families of 2-curves  $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3$  and skeletons  $(\gamma_1, \mathcal{U}_1), (\gamma_2, \mathcal{U}_2), (\gamma_3, \mathcal{U}_3)$  with the following properties:

- $\mathcal{F} = \mathcal{F}_0 \supseteq \mathcal{F}_1 \supseteq \mathcal{F}_2 \supseteq \mathcal{F}_3$ ,
- for  $i \in \{1, 2, 3\}$ ,  $(\gamma_i, \mathcal{U}_i)$  is an  $L(\mathcal{F}_{i-1})$ -skeleton or an  $R(\mathcal{F}_{i-1})$ -skeleton supporting  $\mathcal{F}_i$ ,
- $\chi(\mathcal{F}_1) > f_2(f_2(\beta))$ ,  $\chi(\mathcal{F}_2) > f_2(\beta)$ , and  $\chi(\mathcal{F}_3) > \beta$ .

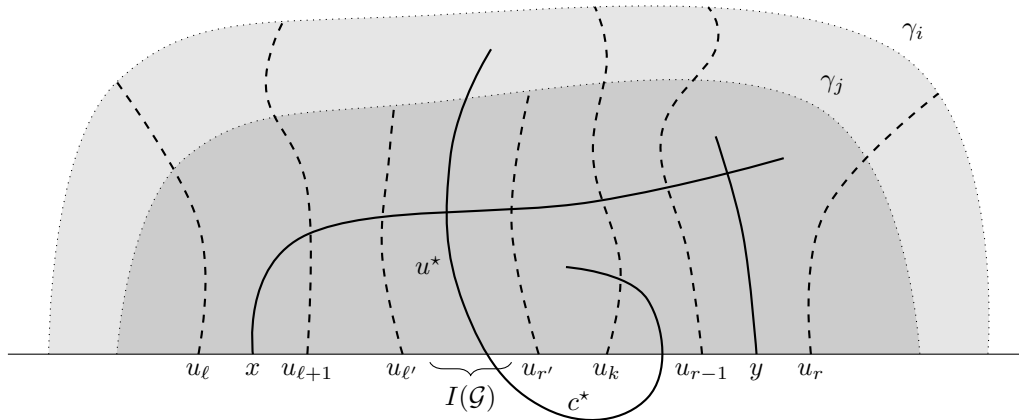


Figure 4 Illustration for the proof of Lemma 19.

There are two indices  $i, j \in \{1, 2, 3\}$  with  $i < j$  such that the skeletons  $(\gamma_i, \mathcal{U}_i)$  and  $(\gamma_j, \mathcal{U}_j)$  are of the same “type”: either an  $L(\mathcal{F}_{i-1})$ -skeleton and an  $L(\mathcal{F}_{j-1})$ -skeleton or an  $R(\mathcal{F}_{i-1})$ -skeleton and an  $R(\mathcal{F}_{j-1})$ -skeleton. Assume for the rest of the proof that  $(\gamma_i, \mathcal{U}_i)$  is an  $L(\mathcal{F}_{i-1})$ -skeleton and  $(\gamma_j, \mathcal{U}_j)$  is an  $L(\mathcal{F}_{j-1})$ -skeleton; the argument for the other case is analogous.

By Lemma 12, since  $\chi(\mathcal{F}_j) \geq \chi(\mathcal{F}_3) > \beta$ , there is a subfamily  $\mathcal{H} \subseteq \mathcal{F}_j$  such that  $\chi(\mathcal{H}) > f(n)$  and  $\chi(\mathcal{F}_j(x, y)) > h(2\xi) + 4\xi + 2$  for any two intersecting 1-curves  $x, y \in L(\mathcal{H}) \cup R(\mathcal{H})$ . Since  $\chi(\mathcal{H}) > f(n)$ , the family  $\mathcal{H}$  contains a chain  $((a_1, b_1), \dots, (a_n, b_n))$  of length  $n$ . Let  $x$  and  $y$  be the 1-curves  $R(a_n)$  and  $L(b_n)$  assigned so that  $x \prec y$ . By the definition of a chain,  $x$  and  $y$  intersect, and therefore  $\chi(\mathcal{F}_j(x, y)) > h(2\xi) + 4\xi + 2$ .

Enumerate the 1-curves in  $\mathcal{U}_i$  as  $u_1, \dots, u_m$  so that  $u_1 \prec \dots \prec u_m$ , where  $m = |\mathcal{U}_i|$ . Assume  $u_1 \prec x \prec y \prec u_m$  for simplicity (adjusting the proof to the general case is straightforward). There are indices  $\ell$  and  $r$  with  $1 \leq \ell < r \leq m$ ,  $u_\ell \prec x \prec u_{\ell+1}$ , and  $u_{r-1} \prec y \prec u_r$ . Let  $\mathcal{F}_j^L = \{c \in \mathcal{F}_j : x \prec L(c) \prec u_{\ell+1}\}$  and  $\mathcal{F}_j^R = \{c \in \mathcal{F}_j : u_{r-1} \prec R(c) \prec y\}$ . It follows that  $\mathcal{F}_j(x, y) \subseteq \mathcal{F}_j^L \cup \mathcal{F}_j(u_{\ell+1}, u_{r-1}) \cup \mathcal{F}_j^R$ .

Since  $\mathcal{F}$  is a  $\xi$ -family, the 2-curves in  $\mathcal{F}_j^L$  that intersect  $u_\ell$  have chromatic number at most  $\xi$ , and so do the 2-curves in  $\mathcal{F}_j^L$  that intersect  $u_{\ell+1}$ . The remaining 2-curves  $c \in \mathcal{F}_j^L$  (intersecting neither  $u_\ell$  nor  $u_{\ell+1}$ ) are pairwise disjoint, because their 1-curves  $L(c)$  are contained in and  $R(c)$  are disjoint from the part of  $\text{int } \gamma_i$  between  $u_\ell$  and  $u_{\ell+1}$ . Thus  $\chi(\mathcal{F}_j^L) \leq 2\xi + 1$ . Similarly,  $\chi(\mathcal{F}_j^R) \leq 2\xi + 1$ . This yields  $\ell + 1 \leq r - 1$  and  $\chi(\mathcal{F}_j(u_{\ell+1}, u_{r-1})) \geq \chi(\mathcal{F}_j(x, y)) - 4\xi - 2 > h(2\xi)$ .

Since  $\mathcal{F}$  is a  $(\xi, h)$ -family, there is a subfamily  $\mathcal{G} \subseteq \mathcal{F}_j(u_{\ell+1}, u_{r-1})$  with  $\chi(\mathcal{G}) > 2\xi$  such that every 2-curve  $c \in \mathcal{F}$  with a basepoint on  $I(\mathcal{G})$  satisfies  $u_{\ell+1} \prec c \prec u_{r-1}$ .

Let  $u_{\ell'}$  be the 1-curve in  $\mathcal{U}_j$  with rightmost basepoint to the left of  $I(\mathcal{G})$ , and let  $u_{r'}$  be the 1-curve in  $\mathcal{U}_j$  with leftmost basepoint to the right of  $I(\mathcal{G})$ . Every 2-curve in  $\mathcal{G}$  must intersect  $u_{\ell'}$ , some 1-curve in  $\mathcal{U}_j(I(\mathcal{G}))$ , or  $u_{r'}$ . Since  $\mathcal{F}$  is a  $\xi$ -family, the 2-curves in  $\mathcal{G}$  that intersect  $u_{\ell'}$  have chromatic number at most  $\xi$ , and so do the 2-curves in  $\mathcal{G}$  that intersect  $u_{r'}$ . Therefore, since  $\chi(\mathcal{G}) > 2\xi$ , some 2-curve in  $\mathcal{G}$  must intersect a 1-curve in  $\mathcal{U}_j(I(\mathcal{G}))$ . In particular, the family  $\mathcal{U}_j(I(\mathcal{G}))$  is non-empty.

Let  $u^* \in \mathcal{U}_j(I(\mathcal{G}))$ . The 1-curve  $u^*$  is a subcurve of  $L(c^*)$  for some 2-curve  $c^* \in \mathcal{F}_{j-1}$ . Since the basepoint of  $L(c^*)$  lies on  $I(\mathcal{G})$ , the property of  $\mathcal{G}$  implies  $u_{\ell+1} \prec c^* \prec u_{r-1}$ . Since  $c^* \in \mathcal{F}_{j-1} \subseteq \mathcal{F}_i$  and  $\mathcal{F}_i$  is supported by  $(\gamma_i, \mathcal{U}_i)$ , the 1-curve  $R(c^*)$  intersects at least one of the 1-curves  $u_{\ell+1}, \dots, u_{r-1}$ , say  $u_k$ . Let  $a_{n+1} = c^*$  and  $b_{n+1}$  be the 2-curve in  $\mathcal{F}_{i-1}$  such that  $u_k$  is a subcurve of  $L(b_{n+1})$ . For  $1 \leq t \leq n$ , the 1-curves  $R(a_t)$  and  $L(b_t)$  intersect and

they are both contained in  $\text{int } \gamma_j$  (because  $a_t, b_t \in \mathcal{H}$ ), the basepoint of  $L(a_{n+1})$  is between the basepoints of  $R(a_t)$  and  $L(b_t)$ , and  $L(a_{n+1})$  intersects  $\gamma_j$  (as it contains  $u^*$ ). Therefore,  $L(a_{n+1})$  intersects all  $R(a_1), \dots, R(a_n)$ . We conclude that  $((a_1, b_1), \dots, (a_{n+1}, b_{n+1}))$  is a chain of length  $n + 1$ . ◀

**Proof of Lemma 15.** Let  $\zeta = f(2k + 1)$ , where  $f$  is the function claimed by Lemma 19 for  $\xi$  and  $h$ . Suppose  $\chi(\mathcal{F}) > \zeta$ . It follows that  $\mathcal{F}$  contains a chain of length  $2k + 1$ . This chain contains a subchain  $((a_1, b_1), \dots, (a_{k+1}, b_{k+1}))$  of pairs of the same “type”:  $L(a_i)$  intersects  $R(a_1), \dots, R(a_{i-1})$  for  $2 \leq i \leq k + 1$  and thus  $\{a_1, \dots, a_{k+1}\}$  is a clique, or  $R(b_i)$  intersects  $L(b_1), \dots, L(b_{i-1})$  for  $2 \leq i \leq k + 1$  and thus  $\{b_1, \dots, b_{k+1}\}$  is a clique. Thus  $\omega(\mathcal{F}) > k$ . ◀

### 3 Proof of Theorem 2

► **Lemma 20** (Fox, Pach, Suk [9, Lemma 3.2]). *For every  $t \in \mathbb{N}$ , there is a constant  $\nu_t > 0$  such that every family of curves  $\mathcal{F}$  any two of which intersect in at most  $t$  points has subfamilies  $\mathcal{F}_1, \dots, \mathcal{F}_d \subseteq \mathcal{F}$  with the following properties:*

- for  $1 \leq i \leq d$ , there is a curve  $c_i \in \mathcal{F}_i$  intersecting all curves in  $\mathcal{F}_i \setminus \{c_i\}$ ,
- for  $1 \leq i < j \leq d$ , every curve in  $\mathcal{F}_i$  is disjoint from every curve in  $\mathcal{F}_j$ ,
- $|\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d| \geq \nu_t |\mathcal{F}| / \log |\mathcal{F}|$ .

**Proof of Theorem 2.** Let  $\mathcal{F}$  be a family of curves obtained from the edges of  $G$  by shortening them slightly so that they do not intersect at the endpoints but all other intersection points are preserved. It follows that  $\omega(\mathcal{F}) \leq k - 1$  (as  $G$  is  $k$ -quasi-planar) and any two curves in  $\mathcal{F}$  intersect in at most  $t$  points. Let  $\nu_t, \mathcal{F}_1, \dots, \mathcal{F}_d$ , and  $c_1, \dots, c_d$  be as claimed by Lemma 20. For  $1 \leq i \leq d$ , since  $\omega(\mathcal{F}_i \setminus \{c_i\}) \leq \omega(\mathcal{F}) - 1 \leq k - 2$ , Theorem 1 yields  $\chi(\mathcal{F}_i \setminus \{c_i\}) \leq f_t(k - 2)$ . Thus  $\chi(\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d) \leq f_t(k - 2) + 1$ . For every color class  $\mathcal{C}$  in a proper coloring of  $\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d$  with  $f_t(k - 2) + 1$  colors, the vertices of  $G$  and the curves in  $\mathcal{C}$  form a planar topological graph, and thus  $|\mathcal{C}| < 3n$ . Thus  $|\mathcal{F}_1 \cup \dots \cup \mathcal{F}_d| < 3(f_t(k - 2) + 1)n$ . This, the third property in Lemma 20, and the fact that  $|\mathcal{F}| < n^2$  yield  $|\mathcal{F}| < 3\nu_t^{-1}(f_t(k - 2) + 1)n \log |\mathcal{F}| < 6\nu_t^{-1}(f_t(k - 2) + 1)n \log n$ . ◀

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